

# AN UPPER BOUND FOR THE NUMBER OF EULERIAN ORIENTATIONS OF A REGULAR GRAPH

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We establish a new upper bound for the number of Eulerian orientations of a regular graph with even degrees.

## 1. Introduction

An *Eulerian orientation* of an undirected graph  $G$  is an orientation of the edges such that the indegree of each vertex is equal to its outdegree. We denote by  $\varepsilon(G)$  the number of Eulerian orientations of  $G$ . Lower and upper bounds for  $\varepsilon(G)$  have been independently established in [1, 3]:

A. Schrijver has shown in [3] that

$$(S1) \quad \varepsilon(G) \cong \left( 2^{-d} \binom{2d}{d} \right)^n$$

where  $G$  is a loopless  $2d$ -regular graph with  $n$  vertices. The author has shown in [1]<sup>1</sup> that

$$(LV1) \quad \varepsilon(G) \cong 2^d \left( \frac{(2d-1)!!}{d!} \right)^{n-1}$$

The bound (LV1) is slightly better than (S1). In fact for given  $d$ , accordingly with Theorem 2 of [3] the ground numbers are equal, namely  $\binom{2d}{d}$ , (and cannot be improved as function of  $d$  only). The ratio is asymptotically equivalent to  $\sqrt{\pi d}$ .

The upper bound in [3] is

$$(S2) \quad \varepsilon(G) \cong \left( \frac{2d}{d} \right)^{n/2}$$

An upper bound for  $\varepsilon(G)$  is given in [1] in the case  $d=2$  as an application of Proposition 4.1 (see below Lemma 2). For a connected, 4-regular, loopless graph

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<sup>1</sup> Errata. Proposition 5.3 in [1] is misstated: the exponent  $n$  should be  $n-1$ . Page 408 line 3 read  $2/(d!)^n$  instead of  $1/2(d!)^n$ .

$G$  of  $n \geq 4$  vertices we have

$$(LV2) \quad \varepsilon(G) \leq 9 \cdot 2^{n-3}$$

The ground number  $\left(\frac{2d}{d}\right)^{0.5}$  in (S2) cannot be improved for general  $d$ -regular graphs [3] (the worst case is given by graphs whose connected components are  $d$  parallel edges). The bound (LV2), which holds for connected graphs, is strictly better  $\left(\frac{9}{8} \cdot 2^n\right)$  as compared to  $(\sqrt{6})^n$ .

Our purpose in the present note is to derive a new upper bound for  $\varepsilon(G)$  in terms of  $d, n$  and the maximum number of edge-disjoint circuits of  $G$ . Applied to loopless connected graphs this bound is significantly sharper than (S2).

## 2. Main theorem

Graphs considered in the sequel may contain loops. In evaluating  $\varepsilon(G)$  we make the convention that there are two ways to direct a loop.

**Theorem 1.** *Let  $G$  be a connected  $2d$ -regular graph,  $d \geq 2$ . We have*

$$(LV2a) \quad \varepsilon(G) \leq 2^k \left( \frac{(2d-1)!!}{d!} \right)^{n - \frac{k-1}{d-1}}$$

where  $n$  is the number of vertices of  $G$  and  $k$  is any number at least equal to the maximal number of pairwise edge-disjoint cycles of  $G$ .

The proof uses an inductive construction of [1]. Let  $x$  be a vertex of  $G$  and  $l$  be the number of loops attached to  $x$ . Let  $T$  be the set of partitions into pairs of the  $2(d-l)$  non-loop edges incident to  $x$ . We associate with each  $t \in T$  a graph  $G_t$  obtained from  $G$  by deleting  $x$  and all incident edges, and introducing for each pair of edges in a same set of  $t$  a new edge joining the endvertices different from  $x$  of these two edges (see [1] for more details). Clearly  $G_t$  is a  $2d$ -regular graph with  $n-1$  vertices, and  $c(G_t) \leq c(G) - l$ , where  $c(G)$  denotes the maximal number of pairwise edge-disjoint cycles of  $G$ . Note that  $|T| = (2d-2l-1)!! = (2d-2l-1)(2d-2l-3)\dots 5.3.1$ .

**Lemma 2.** *We have*

$$\varepsilon(G) = \frac{2^l}{(d-l)!} \sum_{t \in T} \varepsilon(G_t).$$

**Proof.** By counting in two ways the elements of the set  $\{(\emptyset, t) : \emptyset \text{ is an Eulerian orientation of } G \text{ such that for each } t \in T \text{ exactly one of its edges is directed toward } x\}$  we find  $2^l \sum_{t \in T} \varepsilon(G_t) = (d-l)! \varepsilon(G)$ . ■

Set

$$a(n, k) = 2^k \left( \frac{(2d-1)!!}{d!} \right)^{n - \frac{k-1}{d-1}}.$$

**Lemma 3.** *For given  $n$ ,  $a(n, k)$  is an increasing function of  $k$ .*

**Proof.** We have

$$a(n, k) = \left[ 2 \left( \frac{(2d-1)!!}{d!} \right)^{-\frac{1}{d-1}} \right]^k \left( \frac{(2d-1)!!}{d!} \right)^{n+\frac{1}{d-1}}.$$

It suffices to show that

$$2 \left( \frac{(2d-1)!!}{d!} \right)^{-\frac{1}{d-1}} > 1.$$

Raising to the  $(d-1)$ -th power this inequality amounts to

$$\frac{(2d-1)!!}{d!} = \frac{2d-1}{d} \cdot \frac{2d-3}{d-1} \cdots \frac{3}{2} < 2^{d-1}. \blacksquare$$

**Proof of Theorem 1.** In view of Lemma 3 it suffices to establish (LV2a) when  $k=c(G)$ . The proof is by induction on  $n$ .

If  $n=1$  the graph  $G$  consists of  $d$  loops attached to one vertex. We have  $\varepsilon(G)=2^d$ . Theorem 1 with  $k=c(G)=d$  is  $\varepsilon(G) \leq 2^d$ . Suppose  $n \geq 2$ . Let  $x$  be a vertex of  $G$ , not a cutvertex. Let  $l$  be the number of loops attached to  $x$ ,  $0 \leq l \leq d-1$ . By Lemma 2 we have

$$\varepsilon(G) = \frac{2^l}{(d-l)!} \sum_{t \in T} \varepsilon(G_t).$$

There are  $|T|=(2d-2l-1)!!$  graphs  $G_t$ . Each  $G_t$  is connected,  $2d$ -regular with  $n-1$  vertices and  $c(G_t) \leq c(G)-l$ . Hence, using the induction hypothesis and Lemma 3 we get

$$\begin{aligned} \varepsilon(G) &\leq \frac{2^l}{(d-l)!} (2d-2l-1)!! a(n-1, k-l) = \\ &= \frac{(2d-2l-1)!!}{(d-l)!} \cdot \left( \frac{(2d-1)!!}{d!} \right)^{-1+\frac{1}{d-1}} \cdot a(n, k). \end{aligned}$$

To establish Theorem 1 it suffices to show that for  $d \geq 2$  and  $0 \leq l \leq d-1$  we have

$$\frac{(2d-2l-1)!!}{(d-l)!} \cdot \left( \frac{(2d-1)!!}{d!} \right)^{-1+\frac{1}{d-1}} \leq 1.$$

For  $l=d-1$  this inequality holds trivially. Suppose  $0 \leq l \leq d-2$ . The inequality to prove can be rewritten

$$\left( \frac{(2d-2l-1)!!}{(d-l)!} \right)^{\frac{1}{d-l-1}} \leq \left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{d-1}}.$$

Set

$$b(d) = \left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{d-1}}.$$

We show that  $b(d)$  is an increasing function of  $d$  for  $d \geq 2$ , integer. The inequality  $b(d) < b(d+1)$  is

$$\left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{d-1}} < \left( \frac{(2d+1)!!}{(d+1)!} \right)^{\frac{1}{d}}.$$

Raising to the  $d(d-1)$ -th power, we get equivalently

$$\left( \frac{(2d-1)!!}{d!} \right)^d < \left( \frac{(2d+1)!!}{(d+1)!} \right)^{d-1}$$

or, after simplification

$$\frac{(2d-1)!!}{d!} = \frac{2d-1}{d} \cdot \frac{2d-3}{d-1} \cdots \frac{3}{2} < \left( \frac{(2d+1)!!}{(d+1)!} \right)^{d-1}. \quad \blacksquare$$

### 3. Application

**Theorem 4.** *Let  $G$  be a connected  $2d$ -regular graph on  $n$  vertices with  $d \geq 2$  and girth  $\geq g$ . Then*

$$(LV2b) \quad \varepsilon(G) \leq \left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{d-1}} \left[ 2^{\frac{d}{g}} \left( \frac{(2d-1)!!}{d!} \right)^{1 - \frac{1}{g} \cdot \frac{d}{d-1}} \right]^n.$$

**Proof.** Since every cycle of  $G$  is of length  $\geq g$ , we have  $c(G) \leq \frac{dn}{g}$ . Applying

Theorem 1 with  $k = \frac{dn}{g}$ , we get (LV2b).  $\blacksquare$

When  $d \rightarrow \infty$  the ratio of the ground numbers in (S2) and (LV2b) has the asymptotic value

$$2^{\frac{1}{g}} (\sqrt{\pi d})^{\frac{1}{2} - \frac{1}{g}}.$$

A loopless connected graph has girth  $g \geq 2$ . In this case by Theorem 4 we have

$$(LV2b) \quad \varepsilon(G) \leq \left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{d-1}} \left[ 2^{\frac{d}{2}} \left( \frac{(2d-1)!!}{d!} \right)^{\frac{1}{2} \cdot \frac{d-2}{d-1}} \right]^n.$$

As

$$2^d \left( \frac{(2d-1)!!}{d!} \right) = \binom{2d}{d},$$

the ground number in (LV2b) is strictly smaller than the ground number in (S2). The limit of the ratio when  $d \rightarrow \infty$  is  $\sqrt{2}$ . In particular (LV2b) implies (S2) for connected graphs. For  $d=2$  (LV2b) is  $\varepsilon(G) \leq \frac{3}{2} \cdot 2^n$ , slightly weaker than (LV2) (obtained by a finer use of Lemma 2).

**Example.** Applying Theorem 4 with  $g=3$  and  $n=2d+1$  to the complete graph  $K_n$  on  $n$  vertices, we get an upper bound (LV2b) of the asymptotic form

$$4^{d^2} d^{-\frac{2d}{3}(1+o(d))} \quad \text{when } d \rightarrow \infty.$$

In this case (S2) has the form  $4^{d^2} d^{-\frac{d}{2}(1+o(d))}$ , (S1) and (LV1) have the form  $4^{d^2} d^{-d(1+o(d))}$ .

The exact asymptotic value of  $\varepsilon(K_n)$  has recently been obtained by B. McKay [2]. We have

$$\varepsilon(K_n) \cong \left( \frac{2^{n+1}}{\pi n} \right)^{\frac{n-1}{2}} n^{\frac{1}{2}} e^{-\frac{1}{2}}$$

which is of the form  $4^{d^2} d^{-d(1+o(d))}$ .

### References

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